

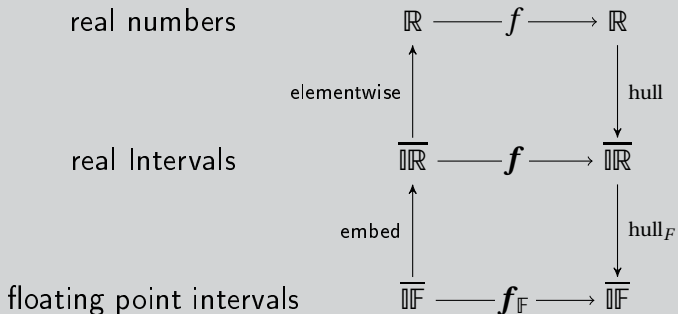
# Computing Interval Power Functions

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## Evaluation



## Definition

Every interval function  $\mathbf{f}$  that contains the range  $f(\mathbf{x})$  is called an interval extension of  $f$

$$x \in \mathbf{x} \Rightarrow f(x) \in \mathbf{f}(\mathbf{x}), \text{ if } f(x) \text{ is defined}$$

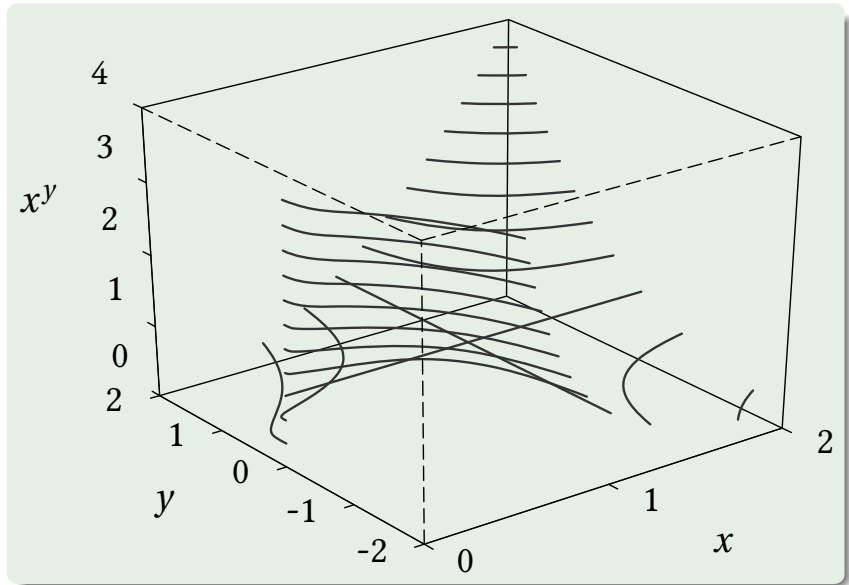
## FTIA

Every interval version of an arithmetic expression is an interval extension of the function defined by that expression.

## Example

Naive interval evaluation of an arithmetic expression.

# General exponential function $x^y$



$\text{pow} : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$

1 positive integral exponents

$$x^n = x \cdot \dots \cdot x$$

2 integral exponents

$$x^{-n} = 1/x^n$$

3 rational exponents

$$x^{m/n} = \sqrt[n]{x^m}$$

4 real exponents

$$x^y = \exp(y \cdot \log x)$$

## Problems

- $0^0?$
- $\sqrt[n]{x^m} \in \mathbb{R}$  if  $x < 0$ ?
- $\exp(y \cdot \log x) \in \mathbb{R}$  if  $x \leq 0$ ?

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$$\text{pow1} \quad \mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log(x))$$

### pow1: positive Version

- Definition only for  $x > 0$
- sufficient for many applications
- IEEE 754 *pow* and *powr*
- mathematical well founded
- differentiable
- restricted, smooth domain and range

$$\text{pow1} \quad \mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

$$\text{powzero} \quad \{0\} \times \mathbb{R}^+ \quad (x, y) \mapsto 0$$



$$\mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

*powzero*       $\{0\} \times \mathbb{R}^+ \quad (x, y) \mapsto 0$

$$\text{pow1} \quad \mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

$$\text{pow2} \quad \{0\} \times \mathbb{R}^+ \quad (x, y) \mapsto 0$$

$$\mathbb{R}^- \times \mathbb{Z} \quad (x, y) \mapsto \begin{cases} \exp(y \cdot \log |x|) & \text{if } y \text{ even} \\ -\exp(y \cdot \log |x|) & \text{if } y \text{ odd} \end{cases}$$

$$\mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

pow2

$$\{0\} \times \mathbb{R}^+ \quad (x, y) \mapsto 0$$

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### pow2: limited Version

- comprises all (commonly agreed) cases of real exponentiation
- restricted, partially discrete domain and range

## complex Version

- uses principal branch of complex logarithm and complex exponential function
- $\exp(y \cdot \log x) \in \mathbb{C}$
- well understood in pure mathematics
- not defined for  $x = 0$ , continuous for  $y > 0$
- as complex function **not suitable**
- But: restriction to real domain yields pow2

$$\text{pow1} \quad \mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

$$\{0\} \times \mathbb{R}^+ \quad (x, y) \mapsto 0$$

$$\text{pow3} \quad \mathbb{R}^- \times \mathbb{Q}_{\text{odd}} \quad (x, \frac{m}{n}) \mapsto \begin{cases} |x|^{m/n} & \text{if } m \text{ even} \\ -|x|^{m/n} & \text{if } m \text{ odd} \end{cases}$$

$\mathbb{Q}_{\text{odd}}$  : odd denominators

$$\mathbb{R}^+ \times \mathbb{R} \quad (x, y) \mapsto \exp(y \cdot \log x)$$

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$\mathbb{Q}_{\text{odd}}$  : odd denominators

### pow3: extended Version

- for  $x < 0$  NOT continuous
- very crisp domain
- contentious application, but may be helpful



$$\mathbf{pow1}(\mathbf{x}, \mathbf{y}) = \text{hull}(\{\text{pow1}(x, y) \mid x \in \mathbf{x} \text{ und } y \in \mathbf{y}\})$$

$$\mathbf{pow2}(\mathbf{x}, \mathbf{y}) = \text{hull}(\{\text{pow2}(x, y) \mid x \in \mathbf{x} \text{ und } y \in \mathbf{y}\})$$

$$\mathbf{pow3}(\mathbf{x}, \mathbf{y}) = \text{hull}(\{\text{pow3}(x, y) \mid x \in \mathbf{x} \text{ und } y \in \mathbf{y}\})$$

$$\text{pow1}(x, y) = \exp(y \cdot \log x)$$

$$\begin{aligned} [\underline{l}, \bar{l}] &= \mathbf{log}_{\mathbb{F}}[\underline{x}, \bar{x}] \\ &= [\nabla \log \underline{x}, \Delta \log \bar{x}], \end{aligned}$$

$$\begin{aligned} [\underline{m}, \bar{m}] &= [\underline{y}, \bar{y}] \bullet_{\mathbb{F}} [\underline{l}, \bar{l}] \\ &= [\min\{\nabla(\underline{y} \cdot \underline{l}), \nabla(\underline{y} \cdot \bar{l}), \nabla(\bar{y} \cdot \underline{l}), \nabla(\bar{y} \cdot \bar{l})\}, \\ &\quad \max\{\Delta(\underline{y} \cdot \underline{l}), \Delta(\underline{y} \cdot \bar{l}), \Delta(\bar{y} \cdot \underline{l}), \Delta(\bar{y} \cdot \bar{l})\}], \end{aligned}$$

$$\begin{aligned} [\underline{z}, \bar{z}] &= \mathbf{exp}_{\mathbb{F}}[\underline{m}, \bar{m}] \\ &= [\nabla \exp \underline{m}, \Delta \exp \bar{m}], \end{aligned}$$

## Goal

Reduce number of operations in floating-point

## Approach

Distinction of cases

- $1 \in \mathbf{x}$ ?
- $0 \in \mathbf{y}$ ?

The value of **pow1**( $[\underline{x}, \bar{x}], [\underline{y}, \bar{y}]$ ) with  $0 < \underline{x}$

	$\bar{x} \leq 1$	$\underline{x} < 1 < \bar{x}$	$1 \leq \underline{x}$
$0 \leq \underline{y}$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$
$\underline{y} < 0 < \bar{y}$	$[\underline{x}^{\underline{y}}, \underline{x}^{\bar{y}}]$	$\text{hull}([\underline{x}^{\underline{y}}, \underline{x}^{\bar{y}}] \cup [\bar{x}^{\underline{y}}, \bar{x}^{\bar{y}}])$	$[\bar{x}^{\underline{y}}, \bar{x}^{\bar{y}}]$
$\bar{y} \leq 0$	$[\bar{x}^{\underline{y}}, \underline{x}^{\underline{y}}]$	$[\bar{x}^{\underline{y}}, \underline{x}^{\underline{y}}]$	$[\bar{x}^{\underline{y}}, \underline{x}^{\underline{y}}]$

## Result

### Speed-up vs. Intlab

- about 45 %
- (Intlab v6; Athlon 64 X2 4850e)

The value of  $\mathbf{pow1}([\underline{x}, \bar{x}], [\underline{y}, \bar{y}])$  with  $0 < \underline{x}$

	$\bar{x} \leq 1$	$\underline{x} < 1 < \bar{x}$	$1 \leq \underline{x}$
$0 \leq \underline{y}$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$	$[\underline{x}^{\underline{y}}, \bar{x}^{\underline{y}}]$
$\underline{y} < 0 < \bar{y}$	$[\underline{x}^{\underline{y}}, \underline{x}^{\bar{y}}]$	$\text{hull}([\underline{x}^{\underline{y}}, \underline{x}^{\bar{y}}] \cup [\bar{x}^{\underline{y}}, \bar{x}^{\bar{y}}])$	$[\bar{x}^{\underline{y}}, \bar{x}^{\bar{y}}]$
$\bar{y} \leq 0$	$[\bar{x}^{\underline{y}}, \underline{x}^{\bar{y}}]$	$[\bar{x}^{\underline{y}}, \underline{x}^{\bar{y}}]$	$[\bar{x}^{\underline{y}}, \underline{x}^{\bar{y}}]$

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## Lemma

Let  $[\underline{x}, \bar{x}], [\underline{y}, \bar{y}] \in \overline{\mathbb{D}}$  where  $\mathbb{D}$  is IEEE-754 binary64 format.  
 Let

$$[\underline{z}, \bar{z}] = \mathbf{exp}_{\mathbb{D}}([\underline{y}, \bar{y}] \bullet_{\mathbb{D}} \mathbf{log}_{\mathbb{D}}[\underline{x}, \bar{x}]).$$

then each of the normal, finite interval boundaries  $\underline{z}$  or  $\bar{z}$  has a worst-case relative error of  $\tilde{\epsilon} = 2^{-41}$  compared to the exact boundary.

improvement needed

- Lauter, Lefevre
- cllibm

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$$\text{pow2} : (\mathbb{R}^+ \times \mathbb{R}) \cup (\{0\} \times \mathbb{R}^+) \cup (\mathbb{R}^- \times \mathbb{Z}) \rightarrow \mathbb{R},$$

$$(x, y) \mapsto \begin{cases} \exp(y \cdot \log x) & \text{if } x \text{ positive,} \\ 0 & \text{if } x \text{ zero,} \\ \exp(y \cdot \log |x|) & \text{if } x \text{ negative and } y \text{ even,} \\ -\exp(y \cdot \log |x|) & \text{if } x \text{ negative and } y \text{ odd.} \end{cases}$$

$$\text{pow2} : (\overline{\mathbb{R}} \times \overline{\mathbb{R}}) \rightarrow \overline{\mathbb{R}},$$

$$(\mathbf{x}, \mathbf{y}) \mapsto \text{hull } \text{pow2}(\mathbf{x}, \mathbf{y})$$



## pow1 $\rightarrow$ pow2

- watch for integral exponents
- negative for odd exponent

## pow1 → pow3

- $\text{pow3}\left(x, \frac{m}{n}\right) = \begin{cases} |x|^{m/n} & \text{if } m \text{ even} \\ -|x|^{m/n} & \text{if } m \text{ odd} \end{cases}$

$(x < 0, n \text{ odd})$

- for negative bases positive and negative powers on dense subsets
- interval extension “wipes out” sign

$x^y$	<b><i>pow1</i></b>	<b><i>pow2</i></b>	<b><i>pow3</i></b>
$[3, 3]^{[2,2]}$	$[9, 9]$	$[9, 9]$	$[9, 9]$
$[-3, -3]^{[2,2]}$	$\emptyset$	$[9, 9]$	$[9, 9]$
$[-3, 2]^{[2,2]}$	$[0, 4]$	$[0, 9]$	$[0, 9]$
$[-3, -3]^{[2,4]}$	$\emptyset$	$[-27, 81]$	$[-81, 81]$
$[-3, 2]^{[-2,3]}$	$[0, +\infty]$	$[-\infty, +\infty]$	$[-\infty, +\infty]$
$[-3, 0]^{[-2,3]}$	$[-\infty, +\infty]$	$[-\infty, +\infty]$	$[-\infty, +\infty]$
$[0, 2]^{[-2,3]}$	$[0, \infty]$	$[0, +\infty]$	$[0, +\infty]$
$[-3, 2]^{[-2,0]}$	$[1/4, +\infty]$	$[-\infty, +\infty]$	$[-\infty, +\infty]$
$[-3, 2]^{[0,3]}$	$[0, 8]$	$[-3, 8]$	$[-27, +27]$
$[-9, -9]^{[1/2, 1/2]}$	$\emptyset$	$\emptyset$	$\emptyset$
$[-8, -8]^{[\nabla(1/3), \Delta(1/3)]}$	$\emptyset$	$\emptyset$	$\supseteq [-2, 2]$

- Discussion of 3 or 4 exponential function(s)
- mathematically founded
- algorithms for all variants but
- extended version as option
- Proof of Concept: Reference implementation  
<http://exp.ln0.de/>
- Marco's talk : reverse mode

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## Simplifying domain $\widetilde{\text{pow3}}$

- dual-valued function  $\widetilde{\text{pow3}} : (x, y) \mapsto \{\pm|x|^y\}$  für  $x < 0$  und alle  $y \in \mathbb{R}$ .
- easy to calculate
- Interval extension using **pow1**
- Difference between **pow3** und  $\widetilde{\text{pow3}}$  only for point-intervals  $\mathbf{y} = [y, y] = \{y\}$

$$\widetilde{\text{pow3}} : (\mathbb{R}^+ \times \mathbb{R}) \cup (\{0\} \times \mathbb{R}^+) \cup (\mathbb{R}^- \times \mathbb{R}) \rightarrow \wp(\mathbb{R})$$

$$(x, y) \mapsto \begin{cases} \{\exp(y \cdot \log x)\} & \text{if } x \text{ positive,} \\ \{0\} & \text{if } x \text{ zero,} \\ \{\pm \exp(y \cdot \log |x|)\} & \text{if } x \text{ negative,} \end{cases}$$

## Lemma

For negative base intervals  $\mathbf{x}$  and exponent intervals  $\mathbf{y}$  it holds

$$\widetilde{\mathbf{pow3}}(\mathbf{x}, \mathbf{y}) = \text{hull}\{\pm\mathbf{pow1}(-\mathbf{x}, \mathbf{y})\} = \\ \text{hull}(-\mathbf{pow1}(-\mathbf{x}, \mathbf{y}) \cup \mathbf{pow1}(-\mathbf{x}, \mathbf{y})).$$

## Lemma

For base intervals  $\mathbf{x}$  and exponent intervals  $\mathbf{y} = [\underline{y}, \bar{y}]$  with  $\underline{y} < \bar{y}$  it holds  $\mathbf{pow3}(\mathbf{x}, \mathbf{y}) = \widetilde{\mathbf{pow3}}(\mathbf{x}, \mathbf{y})$ .

$$\mathbf{pow3}_{\mathbb{F}} : \overline{\mathbb{R}} \times \overline{\mathbb{R}} \rightarrow \overline{\mathbb{R}}$$

$$(x, y) \mapsto$$

Tabelle: The value of  $\mathbf{pow2}([\underline{x}, \bar{x}], \{n\})$  with  $\bar{x} < 0$  and  $n \in \mathbb{Z}$

	$n$ even	$n$ odd
$0 \leq n$	$[\bar{x}^n, \underline{x}^n]$	$[\underline{x}^n, \bar{x}^n]$
$n \leq 0$	$[\underline{x}^n, \bar{x}^n]$	$[\bar{x}^n, \underline{x}^n]$